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# AN INVESTIGATION OF THE RANGE OF A BOOLEAN FUNCTION

by

Norman H. Eggert, Jr.

A thesis submitted in partial fulfillment of the requirements for the degree

of

MASTER OF SCIENCE

in

Mathematics

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Norman H. Eggert, Jr.

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#### INTRODUCTION TO BOOLEAN ALGEBRA

The purpose of this section is to define a boolean algebra and to determine some of the important properties of it.

A boolean algebra is a set B with two binary operations, join and meet, denoted by + and juxtaposition respectively, and a unary operation, complementation, denoted by ', which satisfy the following axioms:

- (1) for all  $a,b \in B$  (that is, for all a,b elements of B) a + b = b + a and ab = ba, (the commutative laws),
- (2) for all  $a,b,c \in B$ , a+bc=(a+b)(a+b) and a(b+c)=ab+ac, (the distributive laws),
- (3) there exists  $0 \in B$  such that for each  $a \in B$ , a + 0 = a, and there exists  $1 \in B$  such that for each  $a \in B$ , a = a,
- (4) for each  $a \in B$ , a + a' = 1 and a a' = 0.

If a + e = a for all a in B then 0 = 0 + e = e + 0 = e, so that there is exactly one element in B which satisfies the first half of axiom 3, namely 0. Similarly there is exactly one element in B which satisfies the second half of axiom 3, namely 1.

The O and l as defined above will be called the distinguished elements.

If in the statement of any of the four axioms join, meet, O, and I are replaced by meet, join, 1, O respectively, the axiom remains unchanged. Thus it follows that if a statement can be proved from the axioms, then the statement with join, meet, O, and I replaced by meet, join, I, and O respectively can also be proved. The two statements are called dual statements.

- (i) a a = a and its dual
- (i!) a + a = a

Let c + x = 1 and c x = 0, then c' = c' 1 = c' (x + c) = c' x + c' c = x c' + x c = x (c' + c) = x 1 = x. Thus axiom 4 defines the operation complementation in the respect that x = c' if and only if c + x = 1 and c x = 0.

The above property is also useful to show two elements are equal.

Since l = a + a' = a' + a and 0 = a a' = a' a it follows from above that:

(ii) (a')' = a

Notice that (ii) is a self-dual statement.

(iii') 0 = b 0

The absorption laws thus follow: a + ab = al + ab = a(l + b) = al = a:

(iv) a + a b = a and its dual

(iv') a (a + b) = a

Also:  $a + b a^{\dagger} = (a + b) (a + a^{\dagger}) = (a + b) \dot{1} = a + b$ . Hence:

(v) a + b = a + b a and its dual

(v') a b = a (b + a')

Next the associative law for join will be established. Using the fact a' (b a) = a' (b a) + a a' = a' (b a + a) = a' a = 0 in the form x' [(y + z) x] = 0 it follows that: (x + y) + z = (x + x' y) + z = [x + (0 + x' y)] + z = (x + (x' [(y + z) x] + x' y)) + z = (x + (x' [(y + z) x + y])) + z = (x + x' [(y + z) x + y]) + z = (x + (y + z) x + y]) + z = (x + (y + z) x + (y + z) y]) + (x z + z) = (x + (y + z)) (x + y) + [x + (y + z)] z = (x + (y + z)) (x + y) + [x + (y + z)] z = (x + (y + z)) [(x + y) + z]. Thus <math>a + (b + c) = (c + b) + a = [c + (b + a)] [(c + b) + a] = (a + (b + c)) [(a + b) + c] = (a + b) c, and the associative laws are proved: (vi) <math>a + (b + c) = (a + b) + c and its dual

(vi) a + (b + c) = (a + b) + c and its dual (vi') a (b c) = (a b) c

Define a + b + c to be the common value of a + (b + c) and (a + b) + c. Also define a b c = a (b c) = (a b) c.

From (iii) and (iii') it follows that l = 0 + 1 and 0 = 10 = 01; hence:

(vii) l = 0'

and its dual

(vii') 0 = 1'

Since (a + b) a' b' = a a' b' + b a' b' = 0 + 0 = 0and (a + b) + a' b' = (a + b + a') (a + b + b') = 1 1 = 1,the De Morgan laws follow:

(viii) (a + b)' = a' b'

and its dual

(viii') (a b)' = a' + b'

Define x < y to mean x + y = y.

If x < y and y < z then x + y = y and y + z = z, thus x + z = x + y + z = y + z = z and x < z. Therefore the relation < is transitive.

If x < y and y < x then x + y = y and y + x = x thus x = y. Hence the relation < is a partial order in a boolean algebra. A partial order is a reflexive, anti-symmetric, transitive relation.

Define x > y to mean x y = y. In a sense > is the dual relation of <.

If x < y then x + y = y and y x = (x + y) x = xthus y > x. Similarly if x > y then y < x.

Since x y = x is a necessary and sufficient condition that x + y = y, either condition will be used for x < y. Furthermore x < y and y > x will be used interchangeably.

By (iii) and (iii'), it follows that 0 < a < 1 for all elements a in B.

Let x < y. Then x + y = y, or by De Morgan's law x'y' = y' and y' < x'. Also a(x + y) = ay or ax + ay = ay thus ax < ay. Furthermore a + x + y = (a + x) + (a + y) = a + y thus a + x < a + y.

If x < y then x y' = (x y) y' = 0, also if x y' = 0 then x y = x y + x y' = x (y + y') = x or x < y. Similarly if x < y then by De Morgan's law on the above result x' + y = 1, and if x' + y = 1 then x < y.

Let  $\mathbf{x_1}, \dots, \mathbf{x_n}$  be variables whose common domain is a subset D of a boolean algebra B. A function f is called a boolean function if the rule for the function,  $f(\mathbf{x_1}, \dots, \mathbf{x_n})$ , can be be built up from the variables  $\mathbf{x_1}, \dots, \mathbf{x_n}$  and elements of B by a finite number of operations meet, join, and complemention. The range of the boolean function f is the set  $\mathbf{x_1}, \dots, \mathbf{x_n}$  and there exist  $\mathbf{x_1}, \dots, \mathbf{x_n} \in \mathbf{D}$  for which  $f(\mathbf{x_1}, \dots, \mathbf{x_n}) = \mathbf{x_1} \cdot (\mathbf{x_1}, \dots,$ 

### Theorem I.1

Every boolean function of one variable in B has a rule of the form f(x) = f(1) x + f(0) x'.

Proof: Since a = a(x + x') = ax + ax', and x = 1x + 0x', the statement is true if f(x) = a or f(x) = x.

If g(x) and h(x) are of the required form, that is if  $g(x) = g(1) \times + g(0) \times !$  and  $h(x) = h(1) \times + h(0) \times !$ , then [g(x)]!, g(x) + h(x), and g(x) h(x) are also of the required form, since:  $[g(x)]! = [g(1) \times + g(0) \times !]! = [g(1)]! + x!$  [g(0)]! + x =  $[g(1)]! [g(0)]! + [g(1)]! \times + [g(0)]! \times ! = [g(1)]! [g(0)]! \times + [g(1)]! [g(0)]! \times ! + [g(1)]! \times + [g(0)]! \times ! = [g(1)]! \times + [g(0)]! \times ! + [g(0)]! \times ! + [g(0)]! \times ! = [g(1)]! \times + [g(0)]! \times ! + [g$ 

Since a boolean function of one variable is a finite number of applications of meet, join, and complementation on x and elements of B, it follows by induction that all boolean functions of one variable are in the form  $f(x) = f(1) \times f(0) \times f(0)$ .

Define, in a boolean algebra B,  $\sum_{(e)_B} f(e_1, \dots, e_n)$  and  $\prod_{(e)_B} f(e_1, \dots, e_n)$  to be the join and meet respectively over all combinations such that either  $e_i = 0$  or  $e_i = 1$ , where 0 and 1 are the distinguished elements of B.

Example: 
$$\sum_{(e)_B} f(e_1, e_2) = f(1,1) + f(1,0) + f(0,1) + f(0,0)$$
  
and  $\prod_{(e)_B} f(e_1, e_2) = f(1,1) f(1,0) f(0,1) f(0,0).$ 

Let e equal 0 or 1, define  $x^e$  to be x if e = 1 and  $x^i$  if e = 0.

It then follows that a boolean function of one variable is in the form f(x) = f(1) x + f(0) x' =

$$\sum_{(e)_{B}} f(e_{1}) x^{e_{1}}.$$

Theorem I.2

If f is a boolean function of n variables then

$$f(x_1, \dots, x_n) = \sum_{(e)_B} f(e_1, \dots, e_n) x_1^{e_1} \dots x_n^{e_n}.$$

Proof: The theorem will be proved by induction. If n=1, then from above the statement is true. Assume that for n=k the statement is true. Then since  $f(x_1, \dots, x_k, x_{k+1})$  can be thought of as a boolean function of one variable  $x_{k+1}$ ,  $f(x_1, \dots, x_k, x_k) = f(x_1, \dots, x_k, x_k)$  by the induction hypothesis,  $f(x_1, \dots, x_k, x_k) = f(x_1, \dots, x_k) = f(x_1, \dots, x_k)$ 

$$\sum_{(e)_{B}} f(e_{1}, \dots, e_{k}, 1) x_{1}^{e_{1}} \dots x_{k}^{e_{k}} x_{k+1} + \sum_{(e)_{B}} f(e_{1}, \dots, e_{k}, 0) x_{1}^{e_{1}} \dots x_{k}^{e_{k}} x_{k+1}^{e_{k}} =$$

$$\sum_{(e)_{B}} f(e_{1}, \dots, e_{k+1}) x_{1}^{e_{1}} \dots x_{k+1}^{e_{k+1}}$$
. Thus the statement

is true for k+l whenever it is true for k, and hence true for all positive integers.

Theorem I.3

$$\sum_{(e)_{B}} a_{1}^{e_{1}} \cdots a_{n}^{e_{n}} = 1, a_{i} B.$$

Proof: This also will be proved by induction.

If n=1 then  $\sum_{(e)_{B}} a_{1}^{e_{1}} = a_{1} + a_{1}' = 1$ . Assume the statement

is true for n=k. Then  $\sum_{(e)} a_1^{e_1} \cdots a_k^{e_k} a_{k+1}^{e_{k+1}} =$ 

$$\left[\sum_{(e)_{B}} a_{1}^{e_{1}} \cdots a_{k}^{e_{k}}\right] a_{k+1} + \left[\sum_{(e)_{B}} a_{1}^{e_{1}} \cdots a_{k}^{e_{k}}\right] a_{k+1}' =$$

$$a_{k+1} + a'_{k+1} = 1.$$

#### DEFINITION OF A RELATIVE BOOLEAN ALGEBRA

This section will define the concept of a relative boolean algebra of a boolean algebra, and give some properties of them.

Let a boolean algebra B and a, b elements of B, be given. Define the set  $S_{a,b}$  by  $S_{a,b} = \{x: x \in B, a < x < b\}$ . Hence  $S_{a,b}$  is the set of all elements x of B where a < x < b.

Note that if a < b then a,b  $\in$  S<sub>a,b</sub>. Let  $x \in$  S<sub>a,b</sub>, by transitivity a < b. Thus S<sub>a,b</sub>  $\neq \varphi$ , the empty set, if and only if a < b. In section I it was shown that for all  $x \in B$ , 0 < x < 1; thus S<sub>0,1</sub> = B.

Let  $S_{a,b} \neq \varphi$ , that is a < b. Since a (b x + a) = a and b (b x + a) = b x + a b = b x + a, a < b x + a < b.

Thus for all  $x \in B$ ,  $b x + a \in S_{a,b}$ . If  $x \in S_{a,b}$ , then a < x < b or b x + a = x + a = x.

Throughout this section it will be assumed that a boolean algebra B has been given and that all elements are elements of B.

#### Theorem II.1

 $S_{a,b}$  is closed under meet and join of B; that is if  $x,y \in S_{a,b}$ , then  $x+y, xy \in S_{a,b}$ .

Proof: Let  $x,y \in S_{a,b}$ . From above a < x,y < b and it follows that (x + y) + b = (x + b) + (y + b) = b, and (x + y) a = x a + y a = a; therefore a < x + y < b. Similarly, x y + b = (x + b) (y + b) = b, x y a = (x a) (y a) = a and a < xy < b. Hence x + y and x y are elements of  $S_{a,b}$ .

If  $S_{a,b} = \phi$  then  $S_{a,b}$  is closed under meet and join.

From above note that  $C(x) = b \ x' + a \in S_{a,b}$  if  $S_{a,b} \neq \phi$ . The element C(x) will be called the relative complement in B with respect to  $S_{a,b}$ , or simply the relative complement if no confusion will result. C(x) is defined, of course, for each element x of B. For convenience,  $x^{\circ}$  or  $x^{+}$  will be used to denote C(x). C(x) will be shown to have properties of complementation in the set  $S_{a,b}$ .

Theorem II.2

 $S_{a,b}$  is a boolean algebra, with distinguished elements a, b, where the meet and join are the same as in B and complementation being the relative complement in B with respect to  $S_{a,b}$ , if and only if  $S_{a,b} \neq \varphi$ .

Proof: Assume  $S_{a,b} \neq \phi$ . From above, the three operations, meet, join, and relative complement, are

closed. Meet and join are commutative and distributive since they are the operations in B. If  $x \in S_{a,b}$ , then x + a = x and x b = x, thus a and b are the distinguished elements of  $S_{a,b}$  as defined by axiom 3. Let  $x^{\circ} = b \ x' + a$ , the relative complement of x. For each s of  $S_{a,b}$ ,  $x x^{\circ} = x (b x' + a) = b x x' + a x = a$  and  $x + x^{\circ} = x + b x' + a = x + b = b$ . Thus axiom 4 is satisfied. Therefore the set  $S_{a,b}$  together with the operations meet, join, and the relative complement in B with respect to  $S_{a,b}$  is a boolean algebra.

If  $S_{a,b}$  together with the three operations form a boolean algebra, then  $S_{a,b} \neq \phi$ .

Since the meet and join in the boolean algebra  $S_{a,b}$  are the same as in B, the partial order, <, is also the same, that is if  $x,y \in S_{a,b}$  and if x < y in B then x < y in  $S_{a,b}$ .

A relative boolean algebra of a boolean algebra B is a subset,  $S_{a,b}$ , of B, a < b, together with the operations meet and join of B and the relative complement in B with respect to  $S_{a,b}$ . From theorem II.2, a relative boolean algebra is a boolean algebra.

### Theorem II.3

Let  $S_{a,b}$  be a relative boolean algebra of B and  $\overline{S}_{c,d}$  be a relative boolean algebra of  $S_{a,b}$ .

Then  $\bar{S}_{c,d}$  is a relative boolean algebra of B.

Frcof: Since  $\overline{S}_{c,d}$  is a relative boolean algebra of  $S_{a,t}$ , a < c < d < b. If  $x \in B$  and c < x < d then  $x \in S_{a,b}$ , thus  $\overline{S}_{c,d} = \left\{x : x \in S_{a,b}, \ c < x < d\right\} = \left\{x : x \notin B, \ c < x < d\right\}$ . The meet and join are the same operations in both  $S_{a,b}$  and  $\overline{S}_{c,d}$ , and in B. Thus it remains to be shown that the relative complement in  $S_{a,b}$  with respect to  $\overline{S}_{c,d}$  is the same as the relative complement in  $S_{a,b}$  with respect to  $\overline{S}_{c,d}$ .

Into  $x^\circ = b \ x' + a$  be the complement in  $S_{a,b}$ . Then the relative complement in  $S_{a,b}$  with respect to  $\overline{S}_{c,d}$  is  $d \ x^\circ + c = d \ (bx' + a) + c =$  d b  $x' + d \ a + c = d \ x' + c$  since a < c < d < b. Thus the complement in  $\overline{S}_{c,d}$  is the relative complement in  $S_{c,d}$  in B with respect to  $S_{c,d}$ .

## Theorem II.4

If  $S_{a,b}$  and  $S_{c,d}$  are relative boolean algebras of B and  $S_{c,d} \subset S_{a,b}$ , [C means is a subset of], then  $S_{c,d}$  is a relative boolean algebra of  $S_{a,b}$ .

Froof: Since  $S_{a,b}$  and  $S_{c,d}$  are relative boolean algebras of B, the meet and join are the same as in B. Thus, since  $S_{a,b} \subset S_{c,d}$ , the meet and join in  $S_{c,d}$  is the same as in  $S_{a,b}$ . Also  $S_{c,d} = \{x: x \in S_{a,b}, x \in S_{a,b}$ 

c < x < d. Hence it remains to show that the complement in  $S_{c,i}$  is the relative complement in  $S_{a,b}$  with respect to  $S_{c,i}$ ; that is  $d(bx^i + a) + c = dx^i + c$ . This is immediate since a < c < d < b.

The next theorem gives a connection between relative boolean algebras in  $\ B_{\bullet}$ 

### Theorem II.5

Let  $S_{a,b}$  be a relative boolean algebra of B and  $x^{\circ} = b x^{\dagger} + a$ . If  $x \in S_{c,d}$ , then  $x^{\circ} \in S_{c,d}$ .

Prcof: Let  $x \in S_{c,d}$ . Then c < d, or by section I, d' < x' < c', and b d' + a < b x' + a < b c' + a. Thus  $1^c < x^c < c^c$  or  $x^c \in S_{c,d}$ .

### Theorem II.6

Let  $S_{a,b}$  be a relative boolean algebra of B and  $x^{\circ} = b \, x' + a$ . If  $S_{h,k}$  is a relative boolean algebra of B and  $x^{+}$  is the complement of x in  $S_{h,k}$  then  $S_{k^{\circ},h^{\circ}}$  is a relative boolean algebra with the complement of x being  $[(x^{\circ})^{+}]^{\circ}$ .

Proof: Since  $S_{h,k} \neq \phi$ , it follows from theorem II.5 that  $S_{k^{\circ},h^{\circ}} \neq \phi$ , and thus  $S_{k^{\circ},h^{\circ}}$  is a relative boolean algebra with the complement of x being  $k^{\circ} x' \cdot h^{\circ}$ . But  $[(x^{\circ})^{+}]^{\circ} = [k (b x' + a) + h]^{\circ} = [k (b x' + a) + h]^{\circ}$ 

b [k (bx' + a) + h]' + a = b [h' (bx' + a) + k'] + a = b h' x' + b h' a + b k' + a =  $(b h' x + a x') + (b k' + a) = h^{\circ} x' + k^{\circ}.$ 

The function f defined by  $f(x) = (((x^+)^\circ)^+)^\circ$ , where 'and + are defined as in theorem II.6, will be seen later to be a homomorphism from  $S_{h,k}$  onto  $S_{k^\circ,h^\circ}$ .

Theorem II.7

 $x \in S_{a,b} \cap S_{c,d}$ , where  $\cap$  is the set intersection, if and only if a + c < x < b d; that is,  $S_{a,b} \cap S_{c,d} = S_{a+c,b} \cap S_{c,d} = S_{a+c,b} \cap S_{c,d}$ 

Proof: If  $x \in S_{a,b} \cap S_{c,d}$ , then x < b and x < d.

Therefore x (b d) = (x b) (x d) = x or x < b d.

Also x > a and x > c, and it follows that x > a + c.

Hence a + c < x < b d.

If a + c < x < b d then a < x < b and c < x < d or  $x \in S_{a,b}$  and  $x \in S_{c,d}$ .

### Theoren II.8

The following are equivalent if  $S_{a,b} \neq \varphi$  and  $S_{c,d} \neq \varphi$ .

- (a)  $S_{a,b} \cap S_{c,d} \neq \varphi$
- (b) a + c < b d
- (c) a < d and c < b
- (d)  $a + c \in S_{a,b} \cap S_{c,d}$
- (e) bd€Sa,b∩Sc,d

- (f)  $S_{a,b} \cap S_{c,d}$  is a relative boolean algebra. Proof: Let  $S = S_{a,b} \cap S_{c,d}$ .
- (a) implies (b): Let  $S \neq \phi$ . Then there exists  $x \in S$ , and by theorem II.7 a + c < x < b d. Hence a + c < b d.
- (b) implies (c): Let a + c < b d. Then 0 = (a + c) (b d)! = (a + c) (b! + d!) = a b! + c b! + a d! + c d! = c b! + a d!. Hence c b! = 0 and a d! = 0, or a < d and c < b.
- (c) implies (d): Let a < d and c < b. Since a < b and c < d it follows that (a + c) + b d = (a + c + b) (a + c + d) = (b + c) (a + d) = b d, or a + c < b d. By theorem II.7,  $a + c \in S$ .
- (d) implies (e): If  $a+c \in S$  then by theorem II.7 a+c < b d and hence b d  $\in S$  by theorem II.7.
- (e) implies (f): Since b d  $\in$  S, S  $\neq$   $\phi$ , and by theorems II.7 and II.2, S is a relative boolean algebra.
- (f) implies (a): Since S is a relative boolean algebra, by theorem II.2, S  $\neq \phi$ .

### BOOLEAN FUNCTIONS

In this section the connection between boolean functions and the sets of the form  $S_{a,b}$  will be shown; namely that the range of a boolean function is a relative boolean algebra.

In this and the following sections f will denote a boolean function of n variables, that is  $f(x_1, \dots, x_n) =$ 

$$\sum_{(e)_{B}} f(e_{1}, \dots, e_{n}) x_{1}^{e_{1}} \dots x_{n}^{e_{n}}.$$

#### Theorem III.1

If f is a function in B such that  $f(a_1, \dots, a_n) = a$  and  $f(b_1, \dots, b_n) = b$ , then there exist  $c_1, \dots, c_n$  such that  $f(c_1, \dots, c_n) = a + b + c$  for any c an element of B, and furthermore  $c_i \in S_{a_i} b_i, a_i + b_i$ 

Proof: Let  $c_i = a_i c + b_i c^i$ . Since  $a_i b_i c_i = a_i b_i$  and  $a_i + b_i + c_i = a_i + b_i$ , it follows that  $a_i b_i < c_i < a_i + b_i$ .

Sime  $c_{i}' = (a_{i} c + b_{i} c')' = a_{i}' c + b_{i}' c'$ it follows that  $c_{i}^{e_{i}} = a_{i}^{e_{i}} c + b_{i}^{e_{i}} c'$ . Thus  $f(c_{1}, \dots, c_{n}) = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n}) (a_{1} c + b_{1} c')^{e_{1}} \dots (a_{n} c + b_{n} c')^{e_{n}} = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n}) (a_{1} c + b_{1} c')^{e_{1}} \dots (a_{n} c + b_{n} c')^{e_{n}} = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n}) (a_{1} c + b_{1} c')^{e_{1}} \dots (a_{n} c + b_{n} c')^{e_{n}} = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n}) (a_{1} c + b_{1} c')^{e_{1}} \dots (a_{n} c + b_{n} c')^{e_{n}} = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n}) (a_{1} c + b_{1} c')^{e_{1}} \dots (a_{n} c + b_{n} c')^{e_{n}} = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n}) (a_{1} c + b_{1} c')^{e_{1}} \dots (a_{n} c + b_{n} c')^{e_{n}} = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n}) (a_{1} c + b_{1} c')^{e_{1}} \dots (a_{n} c + b_{n} c')^{e_{n}} = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n}) (a_{1} c + b_{1} c')^{e_{1}} \dots (a_{n} c + b_{n} c')^{e_{n}} = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n}) (a_{1} c + b_{1} c')^{e_{1}} \dots (a_{n} c + b_{n} c')^{e_{n}} = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n}) (a_{1} c + b_{1} c')^{e_{1}} \dots (a_{n} c + b_{n} c')^{e_{n}} = \sum_{i=1}^{n} f(e_{1}, \dots, e_{n})^{e_{n}} \dots (a_{n} c + b_{n} c')^{e_{n}} \dots (a_{n} c + b_{n}$ 

$$\sum_{(e)_{B}} f(e_{1}, \dots, e_{n}) (a_{1}^{e_{1}} c + b_{1}^{e_{1}} c^{\dagger}) \dots (a_{n}^{e_{n}} c + b_{n}^{e_{n}} c^{\dagger}) =$$

a c + b c'. The elements  $c_1$ , °°°,  $c_n$  exhibited have the properties required in the statement of the theorem, and thus the theorem is proved.

The next two theorems are immediate results of theorem III.1.

Theorem III.2

If  $f(a_1, \dots, a_n) = a$  and  $f(b_1, \dots, b_n) = b$  then there exist  $c_1, \dots, c_n$  and  $d_1, \dots, d_n$  such that  $f(c_1, \dots, c_n) = a + b \text{ and } f(d_1, \dots, d_n) = a b \text{ where}$   $c_i, d_i \in S_{a_i} \ b_i, a_i + b_i.$ 

Proof: In theorem III.l let c be equal to a then to b.

The next theorem follows by induction on theorem III.2.

Theorem III.3

If  $f(p_{li}, \dots, p_{ni}) = p_i$  for  $i=1, \dots, m$ , then

there exist  $r_1, \dots, r_n$  and  $s_1, \dots, s_n$  such that

$$f(r_1, \dots, r_n) = \sum_{i=1}^m p_i \quad \text{and} \quad f(s_1, \dots, s_n) = \prod_{i=1}^m p_i. \quad \text{Furtheremore, } r_j, s_j \in S_{c_j, d_j} \quad \text{where } c_j = \prod_{j=1}^m p_{ji} \quad \text{and } d_j = \sum_{j=1}^m p_{ji}.$$

The following theorems in this section are concerned with the restriction of the domain of f.

### Theoren III.4

Let  $x_i \in S_{a,b} \neq \varphi$  and f a boolean function. Then

$$f(e_1, \dots, e_n) x_1^e \dots x_n^e = f(\eta_1, \dots, \eta_n) x_1^e \dots x_n^e$$

where  $\eta_i = a$  if  $e_i = 0$  and  $\eta_i = b$  if  $e_i = 1$ .

Proof: Let  $x_i \in S_{a,b}$  and  $\eta_i = a$  if  $e_i = 0$ 

and 
$$\eta_i = b$$
 if  $e_i = 1$ . Consider  $\eta_i^p = x_i^i$ , where

p; is either 0 or 1. There are four cases to consider:

(1) if 
$$p_i = 0$$
 and  $e_i = 0$ 

(2) if 
$$p_i = 0$$
 and  $e_i = 1$ 

(3) if 
$$p_i = 1$$
 and  $e_i = 0$ 

(4) if 
$$p_i = 1$$
 and  $e_i = 1$ .

Case (1), since 
$$\eta_{i} = a$$
,  $\eta_{i}^{p} = x_{i}^{p} = a' x_{i}^{p} = x_{i}^{p} = x_{i}^{p}$ .

Dase (2), since 
$$\eta_i = b$$
,  $\eta_i^i = x_i^i = b^i x_i = 0$ .

Tase (3), since 
$$\eta_i = a$$
,  $\eta_i^{p_i} x_i^{e_i} = a x_i^{!} = 0$ .

Case (4), since  $\eta_{i} = b$ ,  $\eta_{i}^{p_{i}} \times_{i}^{e_{i}} = b \times_{i} = x_{i} = x_{i}^{e_{i}}$ .

It then follows that if  $p_i = e_i$  then  $\eta_i^p = x_i^e = e_i$ 

 $x_{i}^{e}$ , and if  $p_{i} \neq e_{i}$  then  $\eta_{i}^{p_{i}} x_{i}^{e} = 0$ . Hence, if for

each i,  $p_i = e_i$ , it follows that  $\eta_1^{p_1} x_1^{e_1} \cdots \eta_n^{p_n} x_n^{e_n} =$ 

 $x_1^e$  ...  $x_n^e$ , otherwise  $\eta_1^p$   $x_1^e$  ...  $\eta_n^p$   $x_n^e$  = 0. Thus

 $f(\eta_1, \dots, \eta_n) x_1^{e_1} \dots x_n^{e_n} =$ 

 $[\sum_{(p)_B} f(p_1, \dots, p_n) \eta_1^{p_1} \dots \eta_n^{p_n}] x_1^{e_1} \dots x_n^{e_n} =$ 

 $\sum_{(p)_B} f(p_1, \dots, p_n) \eta_1^{p_1} x_1^{e_1} \dots \eta_n^{p_n} x_n^{e_n} =$ 

 $f(e_1, \dots, e_n) \times_1^{e_1} \dots \times_n^{e_n}$ , since there is only one combination such that  $p_i = e_i$  for each i.

Let  $c = \prod_{(\eta)} f(\eta_1, \dots, \eta_n)$  and  $d = \sum_{(\eta)} f(\eta_1, \dots, \eta_n)$ ,

where the meet and join extend over all combinations such that either  $\eta_{\text{i}}$  = a or  $\eta_{\text{i}}$  = b.

The next theorems show that if f is restricted to  $S_{a,b} \neq \phi$ , then the range is  $S_{c,d}$ , where c and d are defined above.

Theorem III.5

Let f be a boolean function with a, b, c, and d defined above. Then there exist  $c_1, \dots, c_n$  and  $d_1, \dots, d_n$  such that  $f(c_1, \dots, c_n) = c$  and  $f(d_1, \dots, d_n) = d$ , where  $c_i, d_i \in S_{a,b}$ .

Proof: Application of theorem III.3.

Theorem III.6

Let f be a boolean function with a, b, c, and d defined above. Then there exist  $h_1, \dots, h_n$  such that  $f(h_1, \dots, h_n) = h$  and  $h_i \in S_{a,b}$ .

Theorem III.7

Let f be a boolean function and let

$$c = \prod_{(\eta)_{S_{a},b}} f(\eta_{1}, \dots, \eta_{n}) \quad \text{and} \quad d = \sum_{(\eta)_{S_{a},b}} f(\eta_{1}, \dots, \eta_{n}). \quad \text{If}$$

 $x_i \in S_{a,b}$  then  $f(x_1, \dots, x_n) \in S_{c,d}$ .

Proof: Let  $x_i \in S_{a,b}$ . Recall that

$$c' = \sum_{(e)_B} [f(\eta_1, \dots, \eta_n)]^{\dagger}$$
 and that  $\eta_i = a$  if  $e_i = 0$ 

and  $\eta_i = b$  if  $e_i = 1$ . It then follows by theorem III.4

that: 
$$f(x_1, \dots, x_n) + c' =$$

$$\sum_{(e)_{B}} f(e_{1}, \cdots, e_{n}) x_{1}^{e_{1}} \cdots x_{n}^{e_{n}} + c'$$

$$\sum_{(e)_{B}} f(\eta_{1}, \cdots, \eta_{n}) \times_{1}^{e_{1}} \cdots \times_{n}^{e_{n}} + \sum_{(e)_{B}} [f(\eta_{1}, \cdots, \eta_{n})] =$$

$$\sum_{\substack{(e) \\ B}} x_1^{e_1} \cdots x_n^{e_n} + \sum_{\substack{(e) \\ B}} [f(\eta_1, \cdots, \eta_n)]' = 1 + c' = 1.$$

Therefore  $f(x_1, \dots, x_n) > c$ .

Also it follows that:  $f(x_1, \dots, x_n) + d =$ 

$$\sum_{(e)_{B}} f(e_{1}, \dots, e_{n}) x_{1}^{e_{1}} \dots x_{n}^{e_{n}} + d =$$

$$\sum_{(e)_{B}} f(\eta_{1}, \dots, \eta_{n}) \times_{1}^{e_{1}} \dots \times_{n}^{e_{n}} + \sum_{(e)_{B}} f(\eta_{1}, \dots, \eta_{n}) =$$

$$\sum_{(e)_B} f(\eta_1, \dots, \eta_n) = d. \text{ Hence } f(x_1, \dots, x_n) < d.$$

Thus  $f(x_1, \dots, x_n)$  is an element of  $S_{c,d}$  whenever  $x_i \in S_{a,b}$ . Hence the theorem is proved.

Theorem III.8

The range of a function, f, in B when the domain is restricted to  $S_{a,b} \neq \phi$  is  $S_{c,d}$ , where c=

$$TT_{f(\eta_1, \dots, \eta_n)}$$
 and  $d = \sum_{(\eta)_{S_a, b}} f(\eta_1, \dots, \eta_n)$  where the

join and meet are over all combinations such that either  $\eta_{\text{i}} \, = \, a \ \ \, \text{or} \ \ \, \eta_{\text{i}} \, = \, b \, .$ 

Some immediate results of theorem III.8 are:

#### Theorem III.9

The range of a function, f, in b is Shok where

$$h = \prod_{(e)} f(e_1, \dots, e_n)$$
 and  $k = \sum_{(e)} f(e_1, \dots, e_n)$ .

Theorem III.10

Let f be a boolean function whose domain is B.

The range of f is B if and only if 
$$\mathcal{T}_{(e)_B} f(e_1, \dots, e_n) =$$

0 and 
$$\sum_{(e)_{B}} f(e_{1}, \dots, e_{n}) = 1.$$

Another important theorem is:

Theorem III.11

If  $f(0, \circ \circ , 0) < b$  and  $f(1, \circ \circ , 1) > a$  then

$$f(x_1, \dots, x_n) \in S_{a,b}$$
, whenever  $x_i \in S_{a,b}$ .

Proof: Let  $x_i \in S_{a,b}$  and  $f(0, \dots, 0) < b$  and

$$f(1, \dots, 1) > a$$
. It follows that  $f(x_1, \dots, x_n) =$ 

$$f(1, ..., 1) x_1 ... x_n + f(x_1, ..., x_n) > a.$$

Whenever there is an i such that  $e_i \neq 0$  then

$$b \times_1^{e_1} \cdots \times_n^{e_n} = x_1^{e_1} \cdots x_n^{e_n}$$
. Since  $f(0, \dots, 0) b =$ 

 $f(), \dots, 0)$ , it follows that b  $f(x_1, \dots, x_n) =$ 

$$\sum_{(e)_{R}} f(e_{1}, \dots, e_{n}) \ b \ x_{1}^{1} \cdots x_{n}^{e_{n}} =$$

 $\sum_{(e)_B} f(e_1, \dots, e_n) x_1^{e_1} \dots x_n^{e_n} = f(x_1, \dots, x_n).$  Thus  $f(x_1, \dots, x_n) < b.$ 

Therefore  $f(x_1, \dots, x_n) \in S_{a,b}$ .

#### BOOLEAN FUNCTIONS IN A RELATIVE BOOLEAN ALGEBRA

In this section it will be understood that  $\eta_1=a$  whenever  $e_1=0$  and  $\eta_1=b$  whenever  $e_1=1$ . This section will consider the relationship between functions in B and functions in a relative boolean algebra of B.

Let  $S_{a,b}$  be a relative boolean algebra of B and  $x^{\circ} = b x' + a$ .

A useful theorem is:

Theorem IV.1

If 
$$x_i \in S_{a,b}$$
 then  $b \times_1^{e_1} \cdots \times_n^{e_n} + a = x_1^{\eta_1} \cdots \times_n^{\eta_n}$ , where  $x_i^{i} = x_i$  if  $\eta_i = b$  and  $x_i^{i} = x^{\circ}$  if  $\eta_i = a$ .

Proof: Let  $x_i \in S_{a,b}$ . If  $e_i = 0$  for some i, then it follows that there exists a permutation,  $i_1, i_2, \cdots, i_j, i_{j+1}, \cdots, i_n$ , of the first  $n$  positive integers such that  $e^{ik} = 1$  if  $1 \le k \le j$  and  $e^{ik} = 0$  if  $j+\le k \le n$ . Hence:  $b \times_1^{e_1} \cdots \times_n^{e_n} + a = b \times_{i_1} \cdots \times_{i_j} \times_{i_{j+1}} \cdots \times_{i_n}^{i_{j+1}} + a \times_{i_1} \cdots \times_{i_j}^{e_n} + a = b \times_{i_1} \cdots \times_{i_j} \times_{i_{j+1}} \cdots \times_{i_n}^{i_{j+1}} + a \times_{i_1} \cdots \times_{i_n}^{i_n} = x_1^{n_1} \cdots \times_{i_n}^{n_n}$ .

If  $e_i = 1$  for all i, then  $b \times_1^{e_1} \cdots \times_n^{e_n} = x_1^{e_1} \cdots x_n^{e_n}$ 

Hence in all cases the statement is true.

An alternate statement of theorem IV.1 follows.

Theorem IV.la

If 
$$x_i \in S_{a,b}$$
 then a'b  $x_1$   $\cdots$   $x_n = a' x_1 \cdots x_n$ .

A function, g, of n variables of a relative boolean algebra,  $S_{a,b}$ , of B will have the form  $g(x_1, \cdots, x_n) =$ 

 $\sum_{\substack{(\eta)\\S_a,b}} g(\eta_1,\cdots,\eta_n) \ x_1^{\eta_1} \cdots \ x_n^{\eta_n}, \ \text{where the join extends over}$ 

all combinations such that either  $\eta_i = a$  or  $\eta_i = b$ .

 $x^b = x$  and  $x^a = x^\circ = b x' + a$ , the complement in  $S_{a,b}$ .

The next theorem connects functions in B and functions in  $S_{a,b}$ .

If f is a function in B, then the statement "f is a function in  $S_{a,b}$ " means there exists a g in  $S_{a,b}$  such that  $f(x_1, \dots, x_n) = g(x_1, \dots, x_n)$  whenever  $x_i \in S_{a,b}$ .

Theorem IV.2

If f is a function in B, and  $x_i \in S_{a,b}$  implies that  $f(x_1, \dots, x_n) \in S_{a,b}$ , then f is a function in  $S_{a,b}$  when the domain of the variables is restricted to  $S_{a,b}$ .

Proof: Since  $x_i \in S_{a,b}$  and  $f(x_1, \dots, x_n) \in S_{a,b}$ ; it follows from theorem III.4 and IV.1 that  $f(x_1, \dots, x_n) = S_{a,b}$ 

$$f(x_1, \dots, x_n) + a = \sum_{(e)_B} f(\eta_1, \dots, \eta_n) x_1^{\eta_1} \dots x_n^{\eta_n} + a =$$

$$\sum_{(e)_{B}} f(\eta_{1}, \dots, \eta_{n}) (b x_{1}^{\eta_{1}} \dots x_{n}^{\eta_{n}} + a) + a =$$

$$\sum_{(\eta)} f(\eta_1, \dots, \eta_n) (x_1^{\eta_1} \dots x_n^{\eta_n} + a) + a =$$

$$\sum_{(\eta)} f(\eta_1, \dots, \eta_n) x_1^{\eta_1} \dots x_n^{\eta_n} + a =$$

$$\sum_{(\eta)} f(\eta_1, \cdots, \eta_n) x_1^{\eta_1} \cdots x_n^{\eta_n}.$$

Thus  $f(x_1, \dots, x_n)$  has the form of a function in  $S_{a,b}$  whenever  $x_i \in S_{a,b}$ .

The importance of theorem IV.2 is not obvious from its statement. One result is that if a boolean albebra B is imbedded as a relative boolean algebra in another boolean algebra  $\overline{B}$ , that is, the given algebra

B is a relative boolean algebra of a larger boolean algebra  $\overline{B}$ , no new functions from B into B are formed by rules with coefficients from  $\overline{B}$ .

Next it is of interest to examine the extension of a boolean function, that is, given a function g in  $S_{a,b}$  is there a function f in B such that whenever  $\mathbf{x}_i \in S_{a,b}$  then  $\mathbf{f}(\mathbf{x}_1, \dots, \mathbf{x}_n) = \mathbf{g}(\mathbf{x}_1, \dots, \mathbf{x}_n)$ ? It will be shown that a necessary and sufficient condition that a function g in  $S_{a,b}$  and a function f in B have the property that  $\mathbf{f}(\mathbf{x}_1, \dots, \mathbf{x}_n) = \mathbf{g}(\mathbf{x}_1, \dots, \mathbf{x}_n)$ ,  $\mathbf{x}_i \in S_{a,b}$ , is that the following conditions hold:

(i)  $\mathbf{a}' \mathbf{g}(\mathbf{y}_1, \dots, \mathbf{y}_n) < \mathbf{f}(\mathbf{e}_1, \dots, \mathbf{e}_n) < \mathbf{g}(\mathbf{y}_1, \dots, \mathbf{y}_n) + \mathbf{b}'$ (ii)  $\mathbf{g}(\mathbf{b}, \dots, \mathbf{b}) < \mathbf{f}(\mathbf{1}, \dots, \mathbf{1}) < \mathbf{g}(\mathbf{b}, \dots, \mathbf{b}) + \mathbf{b}'$ (iii)  $\mathbf{a}' \mathbf{g}(\mathbf{a}, \dots, \mathbf{a}) < \mathbf{f}(\mathbf{0}, \dots, \mathbf{0}) < \mathbf{g}(\mathbf{a}, \dots, \mathbf{a})$ .

### Theorem IV.3

Let g be a function in  $S_{a,b}$ , and f a function in B, such that  $x_i \in S_{a,b}$  implies  $g(x_1, \dots, x_n) = f(x_1, \dots, x_n)$ . Then a'  $g(\eta_1, \dots, \eta_n) < f(e_1, \dots, e_n) < g(\eta_1, \dots, \eta_n) + b$ .

Proof: By theorem IV.la, a'  $g(\eta_1, \dots, \eta_n) = a' b g(\eta_1, \dots, \eta_n) = a' b f(\eta_1, \dots, \eta_n) = \sum_{e \in \mathbb{Z}} f(\overline{e}_1, \dots, \overline{e}_n) a' b \eta_1^{e_1} \dots \overline{e}_n^{e_n} = (\overline{e}_1, \dots, \overline{e}_n)$ 

 $\sum_{(e)} f(e_1, \dots, e_n)$  at  $\eta_1^1 \dots \eta_n^n$ , where  $\overline{\eta}_1 = a$  whenever  $e_i = 0$  and  $\overline{\eta}_i = b$  whenever  $\overline{e}_i = 1$ . If  $\overline{\eta}_i \neq \eta_i$ , then  $\eta_i^i = a$  and  $\eta_i^i = 0$ . Also if  $\overline{\eta}_i = \eta_i$  then  $\eta_i^i = b$  and  $e_i = e_i$ . Thus alg $(\eta_1, \dots, \eta_n) =$  $f(e_1, \dots, e_n)$  as b. Thus as  $g(\eta_1, \dots, \eta_n) < f(e_1, \dots, e_n)$ . Again, by theorem IV.1,  $g(\eta_1, \dots, \eta_n) + b^{\dagger} =$  $b g(\eta_1, ..., \eta_n) + a + b! =$  $\sum_{(e)} f(e_1, \dots, e_n) \ b \ \eta_1^{e_1} \cdots \eta_n^{e_n} + a + b' =$  $\sum_{\overline{(e)}_{n}} f(\overline{e}_{1}, \dots, \overline{e}_{n}) \eta_{1}^{\overline{\eta}_{1}} \dots \eta_{n}^{\overline{\eta}_{n}} + a + b', \text{ where } \overline{\eta}_{1} = a$ whenever  $e_{i} = 0$  and  $\eta_{i} = b$  whenever  $e_{i} = 1$ . Since a < b and  $\overline{\eta}_i \neq \eta_i$  implies  $\eta^i = a$ , it follows that whenever i exists such that  $\eta_{i} \neq \overline{\eta}_{i}$  then  $\eta_{1} \cdots \eta_{n} = \overline{\eta}_{n}$ a, and  $f(\overline{e}_1, \dots, \overline{e}_n)$   $\eta_1^{\overline{\eta}_1} \dots \eta_n^{\overline{\eta}_n} + a =$  $f(\overline{e}_1, \cdots, \overline{e}_n)$  a + a = a. Thus, since  $\eta^{\overline{\eta}}_i = b$  whenever  $\eta_{i} = \overline{\eta}_{i}, g(\eta_{1}, \circ \circ , \eta_{n}) + b' = f(e_{1}, \circ \circ , e_{n}) b + a + b' = f($  $f(e_1, \dots, e_n) + a + b!$ . Hence  $g(\eta_1, \dots, \eta_n) + b! >$ f(e<sub>1</sub>, ° ° , e<sub>n</sub>).

Thus the theorem has been established.

Theorem IV.4

Let g be a function in  $S_{a,b}$  and f a function in B, such that  $g(x_1, \dots, x_n) = f(x_1, \dots, x_n)$  whenever  $x_i \in S_{a,b}$ . Then  $f(1, \dots, 1) > g(b, \dots, b)$  and  $f(0, \dots, 0) < g(a, \dots, a)$ .

Proof: By definition of a boolean function  $f(b, \circ \circ ,b) = f(1, \circ \circ ,1) \ b + f(0, \circ \circ ,0) \ b'.$  Since  $f(b, \circ \circ ,b) = g(b, \circ \circ ,b) \in S_{a,b}, \ f(b, \circ \circ ,b) < b.$  Therefore  $0 = b' \ f(b, \circ \circ ,b) = f(0, \circ \circ ,0) \ b'.$  Hence  $g(b, \circ \circ ,b) = f(b, \circ \circ ,b) = f(1, \circ \circ ,1) \ b,$ 

and  $g(b, \dots, b) < f(1, \dots, 1)$ .

By theorem IV.3,  $f(0, \dots, 0) [g(a, \dots, a) + b'] = f(0, \dots, 0)$ . From above,  $f(0, \dots, 0) b' = 0$ . Hence  $f(0, \dots, 0) = f(0, \dots, 0) [g(a, \dots, a) + b'] = f(0, \dots, 0) g(a, \dots, a)$ , or  $f(0, \dots, 0) < g(a, \dots, a)$ .

Theorem IV.25

Let g be a function in  $S_{a,b}$  and f a function in B. A necessary and sufficient condition that f be the same as g whenever the domain of f is restricted to  $S_{a,b}$  is that the following conditions be satisfied:

(i) 
$$a' g(\eta_1, \cdot \cdot \cdot, \eta_n) < f(e_1, \cdot \cdot \cdot, e_n) < g(\eta_1, \cdot \cdot \cdot, \eta_n) + b'$$

(ii) 
$$g(b, \dots, b) < f(1, \dots, 1) < g(b, \dots, b) + b$$

(iii) a 
$$g(a, \circ \circ , a) < f(0, \circ \circ , 0) < g(a, \circ \circ , a)$$
.

Proof: The necessity follows from theorems IV.3 and IV.4.

Assume the conditions hold. Since  $f(1, \dots, 1) > 0$ 

g(b, ..., b) > a and f(0, ..., 0) < g(a, ..., a) < b, it

follows by theorem III.11 that  $f(x_1, \dots, x_n) \in S_{a,b}$ 

whenever  $x_i \in S_{a,b}$ . Thus  $f(x_1, \dots, x_n) = f(x_1, \dots, x_n) + a =$ 

 $\sum_{(e)_{B}} f(e_{1}, \dots, e_{n}) \times_{1}^{e_{1}} \dots \times_{n}^{e_{n}} + a >$ 

 $\sum_{\text{(e)}_{B}} \text{at } g(\eta_1, \cdots, \eta_n) \text{ } x_1^{e_1} \cdots \text{ } x_n^{e_n} + \text{a} =$ 

 $\sum_{(e)_{B}} a' g(\eta_{1}, \cdots, \eta_{n}) b x_{1}^{e_{1}} \cdots x_{n}^{e_{n}} + a.$  By theorem IV.1,

 $f(x_1, \dots, x_n) > \sum_{(\eta)} a_1 g(\eta_1, \dots, \eta_n) x_1^{\eta_1} \dots x_n^{\eta_n} + a =$ 

 $a \in g(x_1, ..., x_n) + a = g(x_1, ..., x_n).$ 

Also  $f(x_1, \dots, x_n) = b f(x_1, \dots, x_n) + a$  whenever

 $x_i \in S_{a,b}$  from above. Thus, by theorem IV.1,  $f(x_1, \dots, x_n) =$ 

 $a + b = \sum_{(e)_{B}} f(e_{1}, \dots, e_{n}) \times_{1}^{e_{1}} \dots \times_{n}^{e_{n}} <$ 

 $a + b \leq [g(\eta_1, \dots, \eta_n) + b] x_1^{e_1} \dots x_n^{e_n} =$ 

 $a + \sum_{(e)_B} g(\eta_1, \dots, \eta_n) [b \times 1] \dots \times n + a] =$ 

 $a + \sum_{(\eta)_{S_{a,b}}} g(\eta_1, \dots, \eta_n) x_1^{\eta_1} \dots x_n^{\eta_n} = a + g(x_1, \dots, x_n).$ 

Thus  $f(x_1, \dots, x_n) < g(x_1, \dots, x_n)$ .

Therefore  $f(x_1, \dots, x_n) = g(x_1, \dots, x_n)$  if  $x_i \in S_{a,b}$ .

#### FUNCTIONS BETWEEN RELATIVE BOOLEAN ALGEBRAS

A homomorphism between boolean algebras  $\, B \,$  and  $\, \overline{B} \,$  is a function, h, of one variable such that:

- (1) h(p + q) = h(p) + h(q)
- (2) h(p q) = h(p) h(q)
- (3) h(p') = [h(p)]'

whenever p and q are elements of B.

If h is a homomorphism then  $h(0) = h(p p^i) = h(p) [h(p)]^i = 0$  and  $h(1) = h(p + p^i) = h(p) + [h(p)]^i = 1$ . Hence the distinguished elements are preserved by the mapping h.

A boolean homomorphism is a homomorphism which can be expressed as a boolean function.

This section deals with functions between relative boolean algebras of B where the functions are boolean functions.

Let  $S_{a,b}$  and  $S_{c,d}$  be relative boolean algebras of B. Let h be a boolean function of one variable in B whose domain is  $S_{a,b}$  and whose range is  $S_{c,d}$ . Let h(x + y) = h(x) + h(y) and h(x y) = h(x) h(y), whenever x and y are elements of  $S_{a,b}$ . For all  $q \in S_{a,b}$ , h(q) = h(q + a) = h(q) + h(a). Since the range of h is  $S_{c,d}$ , if  $p \in S_{c,d}$  then there exists a  $q \in S_{a,b}$  such that h(q) = p. Hence p = p + h(a) for

all p in  $S_{c,d}$ . Since the distinguished elements are unique and h(a) satisfies axiom 3, h(a) = c. Similarly it follows that h(b) = d. Let  $x^{\circ} = b x^{\dagger} + a$  and  $x^{\dagger} = d x^{\dagger} + c$ , the relative complements of  $S_{a,b}$  and  $S_{c,d}$  respectively. Also  $h(x^{\circ}) + h(x) = h(x + x^{\circ}) = h(b) = d$  and  $h(x^{\circ}) h(x) = h(x^{\circ} x) = h(a) = c$ , thus  $h(x^{\circ}) = [h(x)]^{\dagger}$ . Therefore if h is a boolean function with domain  $S_{a,b}$  and range  $S_{c,d}$  and h preserves meet and join, then h is a boolean homomorphism.

If h is a boolean homomorphism with domain  $S_{a,b}$  and h(a)=c, h(b)=d, then by theorem III.8 the range of h is  $S_{c,d}$ .

### Theorem V.1

Let  $S_{a,b}$  and  $S_{c,d}$  be relative boolean algebras of B. The following are equivalent:

- (i) ac = ad and b'c = b'd
- (ii) there exists a boolean homomorphism between  $\mathbf{S}_{\text{a,b}}$  and  $\mathbf{S}_{\text{c,d}}$
- (iii) there exists a boolean function, h, in B whose domain is  $S_{a,b}$  and range is  $S_{c,d}$ .

#### Proof:

(i) implies (ii): Consider the function f discribed by f(x) = d x + c,  $x S_{a,b}$ . Since f(a) = d x + c = a c + c = c and f(b) = db + c = b d + b! c + c = d + c = d, it follows from theorem III.8

that the range of f is  $S_{c,d}$ . Since f(x + y) = d(x + y) + c = (dx + c) + (dx + c) = f(x) + f(y) and f(xy) = dxy + c = (dx + c) (dy + c) = f(x) f(y) the binary operations are preserved, and f is a boolean honomorphism from  $S_{a,b}$  to  $S_{c,d}$ .

(ii) implies (iii): Since all boolean homomorphisms have as a range a relative boolean algebra it follows that there exists a boolean function in B with domain  $s_{a,b}$  and  $s_{c,d}$ .

(iii) implies (i): Let f be a function with domain  $S_{a,b}$  and range  $S_{c,d}$ . Then by theorem III.8 f(a) f(b) = c and f(a) + f(b) = d or c = f(1) a + f(0) b' + f(1) f(0) a' b and d = f(1) b + f(0) a'. Thus c = f(1) a = a [f(1) b + f(0) a'] = a d and c = b' = f(0) b' = b' [f(1) b + f(0) a'] = b' d.

Note in theorem V.1 that if there is any boolean function with domain  $S_{a,b}$  and range  $S_{c,d}$ , then there exists a boolean homomorphism with that domain and range. Furthermore the constants a,b,c, and d completely determine the existence of a homomorphism.

If  $S_{c,d}$  is any relative boolean algebra of B, then since  $S_{0,1} = B$  and 0 c = 0 d and 1! c = 1! d, there is a boolean homomorphism from B to  $S_{c,d}$ , namely, f(x) = d x + c.

Theorem V.2

Let  $S_{a,b}$  and  $S_{c,d}$  be relative boolean algebras of B.  $S_{c,d}$  is a homomorphic image of  $S_{a,b}$ , that is there exists a boolean homomorphism from  $S_{a,b}$  to  $S_{c,d}$ , if and only if there exists a relative boolean algebra,  $S_{k,h}$ , of B with complementation °, such that a° = d and b° = c.

Proof: Let  $S_{c,d}$  be a homomorphic image of  $S_{a,b}$ . Then by the previous theorem, a c=a d and b' c=b' d. Thus da' + c = da' + a c + c = da' + da + c = d + c = d + c = d + c = c. Hence if is the relative complement with respect to  $S_{c,d}$ , then  $a^{\circ}=d$  and  $b^{\circ}=c$ , as was to be shown.

Let  $S_{a,b}$ ,  $S_{c,d}$ , and  $S_{h,k}$  be relative boolean algebras of B. Furthermore let  $a^\circ = k \ a' + h = d$  and  $b^\circ = k \ b' + h = c$ . Consider  $f(x) = x^\circ = k \ x' + h$ . Since f(a) + f(b) = d and  $f(a) \ f(b) = c$  it follows from theorem III.8 that f has a range  $S_{c,d}$  when the domain of f is  $S_{a,b}$ . By theorem V.1,  $S_{c,d}$  is a homomorphic image of  $S_{a,b}$ .

Let  $S_{a,b}$  and  $S_{h,k}$  be relative boolean algebras in B and  $x^{\circ}$  be the relative complement of x with respect to  $S_{h,k}$ . By theorems II.6 and V.2,  $S_{b^{\circ},a^{\circ}}$  is a homomorphic image of  $S_{a,b^{\circ}}$ .

There may be many equations for boolean functions in B which are homomorphisms with domain  $S_{a,b}$  and range  $S_{c,d}$ , but the next theorem shows that if f and g are boolean homomorphisms from  $S_{a,b}$  to  $S_{c,d}$  then f(x) = g(x) whenever  $x \in S_{a,b}$ .

### Theorem V.3

Let  $S_{c,d}$  and  $S_{a,b}$  be relative boolean algebras B. If f is a boolean homomorphism with domain  $S_{a,b}$  and range  $S_{c,d}$  then f(x) = dx + c whenever  $x \in S_{a,b}$ .

Proof: Since f is a homomorphism from  $S_{a,b}$  to  $S_{c,d}$ , f(b) = f(1) b + f(0) b' = d or f(1) b = d b. Also f(a) = f(1) a + f(0) a' = c or f(0) a' = c a'. Let  $x \in S_{a,b}$ . Then f(x) = f(1) x + f(0) x' = f(1) b x + f(0) a' x' = d b x + c a' x' = d x + c x' = d x + c.

If  $S_{a,b}$  and  $S_{c,d}$  are relative boolean algebras of B, and there is a boolean homomorphism from  $S_{a,b}$  to  $S_{c,d}$ , then the function f, f(x) = d + c is a boolean homomorphism, as was shown in the proof of theorem V.1.

Theorem V.3 does not give the form of the boolean function in B, but if B is the domain then there is a unique boolean function which is a homomorphism from

B to  $S_{c,d}$ . Theorem V.4 follows from theorem V.3

Theorem V.4

If B is the domain and  $S_{c,d}$  is the range of a boolean homomorphism f then f(x) = d + c.

Furthermore f is the only boolean homomorphism from B onto  $S_{c,d}$ .

Let  $S_{a,b}$  and  $S_{c,d}$  be relative boolean algebras in B and °, <sup>+</sup> be the complementation operations in  $S_{a,b}$  and  $S_{c,d}$  respectively. If  $S_{c,d}$  is the homomorphic image of  $S_{a,b}$ , then the function f defined by  $f(x) = (x^{\circ})^{+}$ , will be shown to be a boolean homomorphism from  $S_{a,b}$  to  $S_{c,d}$ .  $f(x) = (x^{\circ})^{+} = (x b^{\dagger} + a)^{+} = d(b x^{\dagger} + a)^{\dagger} + c = d a^{\dagger} x + d a x + c = d x + c$ , and therefore f is a boolean homomorphism.

Let  $S_{a,b}$  and  $S_{h,k}$  be relative boolean algebras of B and  $x^{\circ} = b \ x' + a$ ,  $x^{+} = k \ x' + h$ . On page 14 it was stated that  $f(x) = (((x^{+})^{\circ})^{+})^{\circ}$  is a homomorphism; this will now be shown. By theorem II.6,  $S_{k^{\circ},h^{\circ}}$  is a relative boolean algebra with the complement of x being  $((x^{\circ})^{+})^{\circ}$ . Thus, from above,  $f(x) = (((x^{+})^{\circ})^{+})^{\circ}$  is a boolean homomorphism from  $S_{h,k}$  to  $S_{k^{\circ},h^{\circ}}$ .

### RELATIVE BOOLEAN ALGEBRAS AND SUB-BOOLEAN ALGEBRAS

A sub-boolean algebra of a boolean algebra B is a subset of B which is closed under meet, join and complementation of B.

Let S be a sub-boolean algebra of B. Let  $x \in S$  then  $x' \in S$  and  $x + x' = 1 \in S$ , also  $x x' = 0 \in S$ . Thus S is a boolean algebra with distinguished elements 0, 1.

This section deals with the relationship between relative boolean algebras and sub-boolean algebras.

Let SRBA mean "sub-boolean algebra of a relative boolean algebra", and RSBA mean "relative boolean algebra of a sub-boolean algebra ".

It will be shown that if S is a SRBA of B then S is a RSBA of B and if S is a RSBA of B then S is a SRBA of B, where B is a boolean algebra.

Theorem VI.1

If  $\overline{S}$  is a sub-boolean algebra of B and  $\overline{S}_{a,b}$  is a relative boolean algebra of  $\overline{S}$ , then  $\overline{S}_{a,b}$  is a sub-boolean algebra of the relative boolean algebra  $S_{a,b}$  of B.

That is, a RSBA is a SRBA.

Proof: Since  $\overline{S}_{a,b} = \{x:a,b,x\in\overline{S}, \text{ and } a < x < b\}$ , it follows that  $\overline{S}_{a,b}$  is a subset of  $S_{a,b}$ . The comple-

ment in the RSBA,  $\overline{S}_{a,b}$ , is defined by  $b \times l + a$ , the same as in  $S_{a,b}$ . Also the meet and join in both  $\overline{S}_{a,b}$  and  $S_{a,b}$  are the same as in B and thus the same. Hence  $\overline{S}_{a,b}$  is a sub-boolean algebra of  $S_{a,b}$ .

Theorem VI.2

Let  $S_{a,b}$  be a relative boolean algebra of B and S a sub-boolean algebra of  $S_{a,b}$ . Then S is a relative boolean algebra of some sub-boolean algebra of B. That is, a SRBA is a RSBA.

Proof: Let  $\overline{S} = \{x:b \times + a \in S\}$ . If  $x,y \in \overline{S}$  then x + y and x y are elements of  $\overline{S}$  since  $b (x + y) + a = (b \times + a) (b y + a)$  and  $b \times y + a = (b \times + a) (b y + a)$  are elements of S. Let  $x \in \overline{S}$ . Then  $b \times i' + a$  is the complement of  $b \times + a$  with respect to  $S_{a,b}$ . Thus  $b \times i' + a \in S$  and  $x \in \overline{S}$ . Hence  $\overline{S}$  is a subalgebra of B. Note that, since  $b \times a + a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$ , then  $a \times a \in S$  and  $b \times a \in S$  and  $b \times a \in S$ .

Let  $\overline{S}_{a,b} = \{x: x \in \overline{S}, a < x < b\}$ .  $\overline{S}_{a,b}$  is a relative boolean albebra of  $\overline{S}$ . If  $x \in \overline{S}_{a,b} \subset \overline{S}$ , then  $x = b \times a$ , and by definition of  $\overline{S}$ ,  $b \times a \in S$ . Hence  $a \in S$  and  $\overline{S}_{a,b} \subset S$ . If  $a \in S \subset S_{a,b}$ , then  $a < a \times b$  and  $a = a \times b$ . Thus  $a \in \overline{S}_{a,b}$ . Therefore  $a \in S \subset S_{a,b}$ . Since the complement of  $a \in S \subset S_{a,b}$  and  $a \in S \subset S_{a,b}$  and  $a \in S \subset S_{a,b}$  and  $a \in S \subset S_{a,b}$ . Therefore  $a \in S \subset S_{a,b}$  and  $a \in S \subset S_{a,b}$  and  $a \in S \subset S_{a,b}$ . Therefore  $a \in S \subset S_{a,b}$  and  $a \in S \subset S_{a,b}$ . Therefore  $a \in S \subset S_{a,b}$  and  $a \in S \subset S_{a,b}$ . Therefore  $a \in S \subset S_{a,b}$  and  $a \in S \subset S_{a,b}$ . Therefore  $a \in S \subset S_{a,b}$  and  $a \in S \subset S_{a,b}$ .